

Basic Formal Properties of Triangular Norms and Conorms

Adam Grabowski
Institute of Informatics
University of Białystok
Poland

Summary. In the article we present in the Mizar system [1], [8] the catalogue of triangular norms and conorms, used especially in the theory of fuzzy sets [13]. The name *triangular* emphasizes the fact that in the framework of probabilistic metric spaces they generalize triangle inequality [2].

After defining corresponding Mizar mode using four attributes, we introduced the following t-norms:

- minimum t-norm `minnorm` (Def. 6),
- product t-norm `prodnorm` (Def. 8),
- Łukasiewicz t-norm `Lukasiewicz_norm` (Def. 10),
- drastic t-norm `drastic_norm` (Def. 11),
- nilpotent minimum `nilmin_norm` (Def. 12),
- Hamacher product `Hamacher_norm` (Def. 13),

and corresponding t-conorms:

- maximum t-conorm `maxnorm` (Def. 7),
- probabilistic sum `probsum_conorm` (Def. 9),
- bounded sum `BoundedSum_conorm` (Def. 19),
- drastic t-conorm `drastic_conorm` (Def. 14),
- nilpotent maximum `nilmax_conorm` (Def. 18),
- Hamacher t-conorm `Hamacher_conorm` (Def. 17).

Their basic properties and duality are shown; we also proved the predicate of the ordering of norms [10], [9]. It was proven formally that drastic-norm is the pointwise smallest t-norm and `minnorm` is the pointwise largest t-norm (`maxnorm` is the pointwise smallest t-conorm and `drastic-conorm` is the pointwise largest t-conorm).

This work is a continuation of the development of fuzzy sets in Mizar [6] started in [11] and [3]; it could be used to give a variety of more general operations on fuzzy sets. Our formalization is much closer to the set theory used within the Mizar Mathematical Library than the development of rough sets [4], the approach which was chosen allows however for merging both theories [5], [7].

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1. PRELIMINARIES

One can verify that $[0, 1]$ is non empty.

Let us consider elements a, b of $[0, 1]$. Now we state the propositions:

- (1) $\min(a, b) \in [0, 1]$.
- (2) $\max(a, b) \in [0, 1]$.
- (3) $a \cdot b \in [0, 1]$.
- (4) $\max(0, a + b - 1) \in [0, 1]$.
- (5) $\min(a + b, 1) \in [0, 1]$.
- (6) Let us consider elements a, b, c of $[0, 1]$. Then $\max(0, \max(0, a + b - 1) + c - 1) = \max(0, a + \max(0, b + c - 1) - 1)$.
- (7) Let us consider an element a of $[0, 1]$. Then $1 - a \in [0, 1]$.

Let us consider elements a, b of $[0, 1]$. Now we state the propositions:

- (8) $a + b - (a \cdot b) \in [0, 1]$. The theorem is a consequence of (7) and (3).
- (9) $\frac{a \cdot b}{a + b - (a \cdot b)} \in [0, 1]$. The theorem is a consequence of (3) and (8).
- (10) If $\max(a, b) \neq 1$, then $a \neq 1$ and $b \neq 1$.
- (11) Let us consider elements x, y of $[0, 1]$. If $x \cdot y = x + y$, then $x = 0$. The theorem is a consequence of (7).

Let us consider elements a, b of $[0, 1]$. Now we state the propositions:

- (12) $\max(a, b) = 1 - \min(1 - a, 1 - b)$.
- (13) $\min(a + b, 1) = 1 - \max(0, 1 - a + (1 - b) - 1)$.
- (14) $\frac{a + b - (2 \cdot a \cdot b)}{1 - (a \cdot b)} \in [0, 1]$. The theorem is a consequence of (7) and (3).

Let f be a binary operation on $[0, 1]$ and a, b be real numbers. Let us observe that $f(a, b)$ is real.

Now we state the propositions:

- (15) Let us consider real numbers a, b , and a binary operation t on $[0, 1]$. Then $t(a, b) \in [0, 1]$.

- (16) Let us consider a binary operation f on $[0, 1]$, and real numbers a, b . Then $1 - f(1 - a, 1 - b) \in [0, 1]$. The theorem is a consequence of (15) and (7).
- (17) Let us consider real numbers x, y, k . Suppose $k \leq 0$. Then
- (i) $k \cdot \min(x, y) = \max(k \cdot x, k \cdot y)$, and
 - (ii) $k \cdot \max(x, y) = \min(k \cdot x, k \cdot y)$.

2. BASIC EXAMPLE OF A TRIANGULAR NORM AND CONORM: MIN AND MAX

Let A be a real-membered set and f be a binary operation on A . We say that f is monotonic if and only if

(Def. 1) for every elements a, b, c, d of A such that $a \leq c$ and $b \leq d$ holds $f(a, b) \leq f(c, d)$.

We say that f has 1-identity if and only if

(Def. 2) for every element a of A , $f(a, 1) = a$.

We say that f has 1-annihilating if and only if

(Def. 3) for every element a of A , $f(a, 1) = 1$.

We say that f has 0-identity if and only if

(Def. 4) for every element a of A , $f(a, 0) = a$.

We say that f has 0-annihilating if and only if

(Def. 5) for every element a of A , $f(a, 0) = 0$.

The scheme *ExBinOp* deals with a non empty, real-membered set \mathcal{A} and a binary functor \mathcal{F} yielding a set and states that

(Sch. 1) There exists a binary operation f on \mathcal{A} such that for every elements a, b of \mathcal{A} , $f(a, b) = \mathcal{F}(a, b)$

provided

- for every elements a, b of \mathcal{A} , $\mathcal{F}(a, b) \in \mathcal{A}$.

The functor minnorm yielding a binary operation on $[0, 1]$ is defined by

(Def. 6) for every elements a, b of $[0, 1]$, $it(a, b) = \min(a, b)$.

Observe that minnorm is commutative, associative, and monotonic and has 1-identity and there exists a binary operation on $[0, 1]$ which is commutative, associative, and monotonic and has 1-identity.

A t-norm is a commutative, associative, monotonic binary operation on $[0, 1]$ with 1-identity. The functor maxnorm yielding a binary operation on $[0, 1]$ is defined by

(Def. 7) for every elements a, b of $[0, 1]$, $it(a, b) = \max(a, b)$.

One can verify that maxnorm is commutative, associative, and monotonic and has 0-identity and there exists a binary operation on $[0, 1]$ which is commutative, associative, and monotonic and has 0-identity.

A t-conorm is a commutative, associative, monotonic binary operation on $[0, 1]$ with 0-identity. Now we state the propositions:

(18) Let us consider a commutative, monotonic binary operation t on $[0, 1]$ with 1-identity, and an element a of $[0, 1]$. Then $t(a, 0) = 0$. The theorem is a consequence of (15).

(19) Let us consider a commutative, monotonic binary operation t on $[0, 1]$ with 0-identity, and an element a of $[0, 1]$. Then $t(a, 1) = 1$. The theorem is a consequence of (15).

Let us note that every commutative, monotonic binary operation on $[0, 1]$ with 1-identity has 0-annihilating and every commutative, monotonic binary operation on $[0, 1]$ with 0-identity has 1-annihilating.

3. FURTHER EXAMPLES OF TRIANGULAR NORMS

The functor prodnorm yielding a binary operation on $[0, 1]$ is defined by

(Def. 8) for every elements a, b of $[0, 1]$, $it(a, b) = a \cdot b$.

Let us observe that prodnorm is commutative, associative, and monotonic and has 1-identity.

The functor probsum-conorm yielding a binary operation on $[0, 1]$ is defined by

(Def. 9) for every elements a, b of $[0, 1]$, $it(a, b) = a + b - (a \cdot b)$.

The functor Lukasiewicz-norm yielding a binary operation on $[0, 1]$ is defined by

(Def. 10) for every elements a, b of $[0, 1]$, $it(a, b) = \max(0, a + b - 1)$.

One can check that Lukasiewicz-norm is commutative, associative, and monotonic and has 1-identity.

The functor drastic-norm yielding a binary operation on $[0, 1]$ is defined by

(Def. 11) for every elements a, b of $[0, 1]$, if $\max(a, b) = 1$, then $it(a, b) = \min(a, b)$ and if $\max(a, b) \neq 1$, then $it(a, b) = 0$.

Now we state the proposition:

(20) Let us consider elements a, b of $[0, 1]$. Then

- (i) if $a = 1$, then $(\text{drastic-norm})(a, b) = b$, and
- (ii) if $b = 1$, then $(\text{drastic-norm})(a, b) = a$, and

(iii) if $a \neq 1$ and $b \neq 1$, then $(\text{drastic-norm})(a, b) = 0$.

Note that drastic-norm is commutative, associative, and monotonic and has 1-identity.

The functor nilmin-norm yielding a binary operation on $[0, 1]$ is defined by (Def. 12) for every elements a, b of $[0, 1]$, if $a + b > 1$, then $it(a, b) = \min(a, b)$ and if $a + b \leq 1$, then $it(a, b) = 0$.

Observe that nilmin-norm is commutative, associative, and monotonic and has 1-identity.

The functor Hamacher-norm yielding a binary operation on $[0, 1]$ is defined by

(Def. 13) for every elements a, b of $[0, 1]$, $it(a, b) = \frac{a \cdot b}{a + b - (a \cdot b)}$.

One can verify that Hamacher-norm is commutative, associative, and monotonic and has 1-identity.

4. BASIC TRIANGULAR CONORMS

The functor drastic-conorm yielding a binary operation on $[0, 1]$ is defined by

(Def. 14) for every elements a, b of $[0, 1]$, if $\min(a, b) = 0$, then $it(a, b) = \max(a, b)$ and if $\min(a, b) \neq 0$, then $it(a, b) = 1$.

5. TRANSLATING BETWEEN TRIANGULAR NORMS AND CONORMS

Let t be a binary operation on $[0, 1]$. The functor conorm t yielding a binary operation on $[0, 1]$ is defined by

(Def. 15) for every elements a, b of $[0, 1]$, $it(a, b) = 1 - t(1 - a, 1 - b)$.

Let t be a t-norm. Let us observe that conorm t is monotonic, commutative, and associative and has 0-identity.

Now we state the propositions:

(21) $\text{maxnorm} = \text{conorm minnorm}$.

PROOF: For every elements a, b of $[0, 1]$, $(\text{maxnorm})(a, b) = 1 - (\text{minnorm})(1 - a, 1 - b)$ by (7), (17), [12, (42)]. \square

(22) Let us consider a binary operation t on $[0, 1]$. Then conorm conorm $t = t$. The theorem is a consequence of (7).

6. THE ORDERING OF TRIANGULAR NORMS (AND CONORMS)

Let f_1, f_2 be binary operations on $[0, 1]$. We say that $f_1 \leq f_2$ if and only if
(Def. 16) for every elements a, b of $[0, 1]$, $f_1(a, b) \leq f_2(a, b)$.

Let us consider a t-norm t . Now we state the propositions:

(23) drastic-norm $\leq t$. The theorem is a consequence of (20).

(24) $t \leq$ minnorm.

Now we state the proposition:

(25) Let us consider t-norms t_1, t_2 . If $t_1 \leq t_2$, then conorm $t_2 \leq$ conorm t_1 .
The theorem is a consequence of (7).

7. TRIANGULAR CONORMS GENERATED FROM T-NORMS

The functor Hamacher-conorm yielding a binary operation on $[0, 1]$ is defined
by

(Def. 17) for every elements a, b of $[0, 1]$, if $a = b = 1$, then $it(a, b) = 1$ and if
 $a \neq 1$ or $b \neq 1$, then $it(a, b) = \frac{a+b-(2 \cdot a \cdot b)}{1-(a \cdot b)}$.

Now we state the proposition:

(26) conorm Hamacher-norm = Hamacher-conorm. The theorem is a consequence of (7).

Let us note that Hamacher-conorm is commutative, associative, and monotonic and has 0-identity.

Now we state the propositions:

(27) conorm drastic-norm = drastic-conorm. The theorem is a consequence of (7).

(28) conorm prodnorm = probsum-conorm. The theorem is a consequence of (7).

One can check that probsum-conorm is commutative, associative, and monotonic and has 0-identity.

The functor nilmax-conorm yielding a binary operation on $[0, 1]$ is defined
by

(Def. 18) for every elements a, b of $[0, 1]$, if $a + b < 1$, then $it(a, b) = \max(a, b)$ and
if $a + b \geq 1$, then $it(a, b) = 1$.

Now we state the proposition:

(29) conorm nilmin-norm = nilmax-conorm. The theorem is a consequence of (7) and (12).

Let us note that nilmax-conorm is commutative, associative, and monotonic and has 0-identity.

The functor BoundedSum-conorm yielding a binary operation on $[0, 1]$ is defined by

(Def. 19) for every elements a, b of $[0, 1]$, $it(a, b) = \min(a + b, 1)$.

Now we state the proposition:

(30) conorm Lukasiewicz-norm = BoundedSum-conorm. The theorem is a consequence of (7) and (13).

One can check that BoundedSum-conorm is commutative, associative, and monotonic and has 0-identity.

Let us consider a t-conorm t . Now we state the propositions:

(31) maxnorm $\leq t$.

(32) $t \leq$ drastic-conorm.

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