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Pseudo-Canonical Formulae are Classical

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Summary. An original result about Hilbert Positive Propositional Calculus introduced in [11] is proven. That is, it is shown that the pseudo-canonical formulae of that calculus (and hence also the canonical ones, see [17]) are a subset of the classical tautologies.

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The notation and terminology used in this paper have been introduced in the following articles: [13], [1], [14], [10], [9], [15], [3], [4], [5], [6], [11], [16], [17], [2], [7], [18], [20], [22], [21], [12], [19], and [8].

1. Preliminaries about Injectivity, Involutiveness, Fixed Points

From now on a, b, c, x, y, z, A, B, C, X, Y denote sets, f, g denote functions, V denotes a SetValuation, P denotes a permutation of V, p, q, r, s denote elements of HP-WFF, and n denotes an element of \mathbb{N} .

Let us consider X and Y. Let f be a relation between X and Y. Note that $id_X \cdot f$ reduces to f and $f \cdot id_Y$ reduces to f.

Now we state the proposition:

(1) Let us consider one-to-one functions f, g. If $f^{-1} = g^{-1}$, then f = g.

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One can verify that there exists a function which is involutive.

Let us consider A. Let us observe that there exists a permutation of A which is involutive.

Now we state the propositions:

- (2) Let us consider an involutive function f. Suppose rng $f \subseteq \text{dom } f$. Then $f \cdot f = \text{id}_{\text{dom } f}$.
- (3) Let us consider a function f. If $f \cdot f = \mathrm{id}_{\mathrm{dom}\, f}$, then f is involutive.
- (4) Let us consider an involutive function f from A into A. Then $f \cdot f = id_A$. The theorem is a consequence of (2).
- (5) Let us consider a function f from A into A. If $f \cdot f = id_A$, then f is involutive. The theorem is a consequence of (3).

Observe that every function which is involutive is also one-to-one.

Let us consider A. Let f be an involutive permutation of A. One can verify that f^{-1} is involutive.

Let n be a non zero natural number. Observe that $[0 \longmapsto n, n \longmapsto 0]$ is without fixpoints.

Let z be a zero natural number. Note that fixpoints $[z \longmapsto n, n \longmapsto z]$ is empty.

Let X be a non empty set. Observe that there exists a permutation of X which is non empty and involutive.

Let us consider A and B. Let f be an involutive function from A into A and g be an involutive function from B into B. Observe that $f \times g$ is involutive.

Let A, B be non empty sets, f be an involutive permutation of A, and g be an involutive permutation of B. Observe that $f \Rightarrow g$ is involutive.

2. Facts about Perm's Fixed Points

Now we state the propositions:

- (6) If x is a fixpoint of $\operatorname{Perm}(P,q)$, then $\operatorname{SetVal}(V,p) \longmapsto x$ is a fixpoint of $\operatorname{Perm}(P,p\Rightarrow q)$.
- (7) If $\operatorname{Perm}(P,q)$ has fixpoints, then $\operatorname{Perm}(P,p\Rightarrow q)$ has fixpoints. The theorem is a consequence of (6).
- (8) If Perm(P, p) has fixpoints and Perm(P, q) is without fixpoints, then $Perm(P, p \Rightarrow q)$ is without fixpoints.

3. Axiom of Choice in Functional Form via the Fraenkel Operator

Let X be a set. The functor Choice On X yielding a set is defined by the term

(Def. 1) $\{\langle x, \text{ the element of } x \rangle, \text{ where } x \text{ is an element of } X \setminus \{\emptyset\} : x \in X \setminus \{\emptyset\}\}.$

One can check that $\operatorname{ChoiceOn} X$ is relation-like and function-like.

Let us consider f. The functor FieldCover f yielding a set is defined by the term

(Def. 2) $\{\{x, f(x)\}\}$, where x is an element of dom $f: x \in \text{dom } f\}$.

The functor SomePoints f yielding a set is defined by the term

(Def. 3) field $f \setminus \operatorname{rng} \text{ChoiceOn FieldCover } f$.

The functor OtherPoints f yielding a set is defined by the term

(Def. 4) (field $f \setminus \text{fixpoints } f$) \ SomePoints f.

Let us consider g. Let us observe that OtherPoints $g \cap SomePoints g$ is empty.

4. Building a Suitable Set Valuation and a Suitable Permutation of It

Let us consider x. The functor ToHilb(x) yielding a set is defined by the term

(Def. 5)
$$(\mathrm{id}_1 + \cdot (1 \times \emptyset^x) \cdot (\emptyset^x \times \{1\})) + \cdot (\{1\} \times \emptyset^x) \cdot (\emptyset^x \times \{0\}).$$

Note that ToHilb(x) is function-like and relation-like.

Now we state the propositions:

- (9) If $x \neq \emptyset$, then ToHilb $(x) = id_1$.
- (10) ToHilb(\emptyset) = $[0 \longmapsto 1, 1 \longmapsto 0]$.

Let v be a function. The functor ToHilbPerm(v) yielding a set is defined by the term

(Def. 6) the set of all (n, ToHilb(v(n))) where n is an element of N.

The functor ToHilbVal(v) yielding a set is defined by the term

(Def. 7) the set of all $\langle n, \text{dom ToHilb}(v(n)) \rangle$ where n is an element of N.

One can check that ToHilbVal(v) is function-like and relation-like and ToHilb-Perm(v) is function-like and relation-like and ToHilbVal(v) is \mathbb{N} -defined and ToHilbVal(v) is total and ToHilbPerm(v) is \mathbb{N} -defined and ToHilbPerm(v) is total.

One can verify that ToHilbVal(v) is non-empty.

Let us consider x. Let us note that ToHilb(x) is symmetric.

Let v be a function. Observe that the functor ToHilbPerm(v) yields a permutation of ToHilbVal(v).

A set valuation is a many sorted set indexed by \mathbb{N} . From now on v denotes a set valuation.

Let us consider p and v. Note that Perm(ToHilbPerm(v), p) is involutive.

5. Classical Semantics via SetVal₀, an Extension of SetVal

Let V be a set valuation. The functor $\operatorname{SetVal}_0 V$ yielding a many sorted set indexed by HP-WFF is defined by

- (Def. 8) (i) it(VERUM) = 1, and
 - (ii) for every n, it(prop n) = V(n), and
 - (iii) for every p and q, $it(p \land q) = it(p) \times it(q)$ and $it(p \Rightarrow q) = (it(q))^{it(p)}$.

Let us consider v and p. The functor $SetVal_0(v, p)$ yielding a set is defined by the term

(Def. 9) (SetVal₀ v)(p).

We say that p is classical if and only if

(Def. 10) SetVal₀ $(v, p) \neq \emptyset$.

One can check that every element of HP-WFF which is pseudo-canonical is also classical.

Let us consider v. Let p be a classical element of HP-WFF. Note that $\operatorname{SetVal}_0(v,p)$ is non empty.

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